

Generalized Interval Embeddings

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Let $\mathbb{N} = \{1, 2, 3, \dots\}$ and let \mathbb{N}^* denote the set of all words over \mathbb{N} . Let ϵ denote the empty word. Given words u and v in \mathbb{N}^* , we say that u is a factor of v if there are words w_1 and w_2 such that $v = w_1uw_2$. In such a situation, we say u is a suffix of v if $w_2 = \epsilon$. Given $u = u_1u_2\dots u_\ell \in \mathbb{N}^*$, we define the *norm* of u to be $\Sigma u = u_1 + u_2 + \dots + u_\ell$ and we define the *length* of u to be $|u| = \ell$. We then allow x and t to be commuting variables and we define the *weight* of u to be $wt(u) = x^{\Sigma u}t^{|u|}$.

Given any poset $\mathcal{P} = (\mathbb{N}, \leq_{\mathcal{P}})$ and $m, n \in \mathbb{N}$, we let $I_{m, \infty}^{\mathcal{P}} = \{k \in \mathbb{N} : m \leq_{\mathcal{P}} k\}$ and $I_{m, n}^{\mathcal{P}} = \{n \in \mathbb{N} : m \leq_{\mathcal{P}} k \leq_{\mathcal{P}} n\}$. Given any words $u = u_1\dots u_k$ and $w = w_1w_2\dots w_\ell$ in \mathbb{N}^* , we say that u *embeds into w relative to \mathcal{P}* , written $u \leq_{\mathcal{P}} w$ if there is a factor $w' = w'_1w'_2\dots w'_k$ of w such that $w'_i \in I_{u_i, \infty}^{\mathcal{P}}$ for every $1 \leq i \leq k$. We define $S^{\mathcal{P}}(u)$ to be the set of all words w that embed u such that the only embedding of u into w occurs at the right end of w , and we set

$$S^{\mathcal{P}}(u, x, t) = \sum_{w \in S^{\mathcal{P}}(u)} wt(w).$$

Given $u, v \in \mathbb{N}^*$, u and v are \mathcal{P} -Wilf equivalent, written as $u \sim_{\mathcal{P}} v$, if $S^{\mathcal{P}}(u, x, t) = S^{\mathcal{P}}(v, x, t)$. Kitaev, Liese, Rempel and Sagan [1] studied various properties of \mathcal{P} -Wilf Equivalence where \mathcal{P} is the standard order on \mathbb{N} and Langley, Liese, and Rempel [2] studied various properties of \mathcal{P}_k -Wilf equivalence where $\mathcal{P}_k = (\mathbb{N}, \leq_k)$ and $i \leq_k j$ if and only if $i \equiv j \pmod k$ and $i < j$.

We study a generalization \mathcal{P} -Wilf equivalence based on intervals. That is, suppose that we are given a poset $\mathcal{P} = (\mathbb{N}, \leq_{\mathcal{P}})$ and a sequence \vec{U} of intervals $(\{I_{m_1, n_1}^{\mathcal{P}}, I_{m_2, n_2}^{\mathcal{P}}, \dots, I_{m_k, n_k}^{\mathcal{P}}\})$ where either $m_i \leq_{\mathcal{P}} n_i$ and $m_i, n_i \in \mathbb{N}$ or $m_i \in \mathbb{N}$ and $n_i = \infty$. Then we say that w has an *interval-embedding of \vec{U} into w relative to \mathcal{P}* , denoted $\vec{U} \leq_{\mathcal{P}} w$ if there is a factor $w' = w'_1w'_2\dots w'_k$ of w such that $w'_i \in I_{m_i, n_i}^{\mathcal{P}}$ for every $1 \leq i \leq k$. We then define $S^{\mathcal{P}}(\vec{U})$ to be the set of all words $w = w_1\dots w_n \in \mathbb{N}^*$ such that $n \geq k$, there is an interval embedding of \vec{U} into the suffix of w of length k , and there is no interval embedding of \vec{U} into $w_1\dots w_{n-1}$. We set

$$S^{\mathcal{P}}(\vec{U}, x, t) = \sum_{w \in S^{\mathcal{P}}(\vec{U})} wt(w),$$

and given two sequences \vec{U} and \vec{V} of intervals of \mathcal{P} , we say that \vec{U} is \mathcal{P} -Wilf equivalent to \vec{V} , written as $\vec{U} \sim_{\mathcal{P}} \vec{V}$, if $S^{\mathcal{P}}(\vec{U}, x, t) = S^{\mathcal{P}}(\vec{V}, x, t)$.

We show that under mild assumptions on \mathcal{P} , $S^{\mathcal{P}}(\vec{U})$ is accepted by a finite automaton and, hence, $S^{\mathcal{P}}(\vec{U}, x, t)$ is a rational function. We compute $S^{\mathcal{P}}(\vec{U}, x, t)$ for various special cases of \vec{U} and use these computations to establish various non-trivial Wilf-equivalences in this setting.

[1] S. Kitaev, J. Liese, J. Rempel, and B.E. Sagan, Rationality of generalized containments in words and Wilf equivalence, *Electron. J. Combin.*, **16(2)** (2009), R22

[2] T. Langley, J. Liese, and J. Rempel, Generating functions for Wilf equivalence under the generalized factor order, *J. Integer Seq.*, **14** (2011), 11.4.2

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