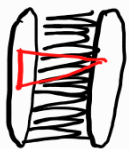


Transversal Embedding (with Y. Cheng)

Scottish Combinatorics
Meeting May '23

Graph embedding: when does a host graph G contain a given subgraph H ?



Computationally difficult - seek (easy to check) sufficient conditions.

Transversal embedding: $s \geq e(H)$ graphs on the same vertex set



Is there a copy of H with ≤ 1 edge from each graph? Copy is **transversal** or **rainbow**.

Generalisation: if every G_i is the same then recover uncoloured problem.

Examples

Edges vs. triangles

not color-blind Mantel 1907: Every n -vx graph with $> \lfloor \frac{n^2}{4} \rfloor$ edges contains a triangle



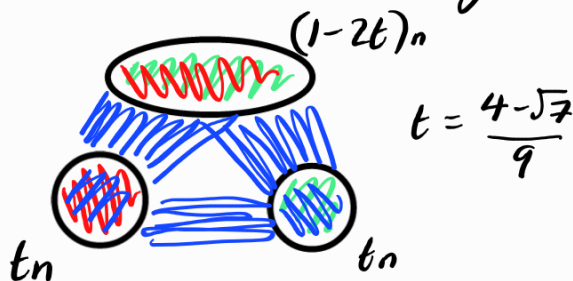
$$\lfloor \frac{n}{2} \rfloor \quad \lfloor \frac{n}{2} \rfloor$$

Aharoni-DeVos-de la Maza-Montejano-Sámal '20

G_1, G_2, G_3 n -vx graphs,

$$e(G_i) > \left(\frac{26 - 2\sqrt{7}}{81} \right) n^2 \approx 0.2557 n^2$$

contain a rainbow triangle



Min degree vs Hamilton cycle

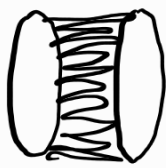
color-blind Dirac 1952: Every n -vx graph with min degree $\geq \frac{n}{2}$ contains a Hamilton cycle.

Joo-Kim '18:

$G_1, \dots, G_n$ n -vx graphs,

$$\delta(G_i) \geq \frac{n}{2} \quad \text{for all } i$$

contain a rainbow Hamilton cycle.



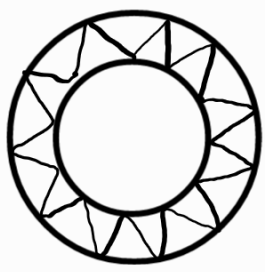
(Joos-Kim) Which embedding problems are colour-blind?

Min degree vs spanning H

approximately colour-blind: $+ \epsilon n$ degree

$\triangle \triangle \triangle \dots \triangle$ Montgomery-Müyesser-Pehova '21

$o\left(\frac{n}{\log n}\right)$



Gupta-Hamann-Müyesser-Parczyk-Szeglia '22+

$\leftarrow o(n) \rightarrow$
bounded degree
small bandwidth

Chakraborti-Im-Kim-Lin '23+

Proof techniques (MMP)

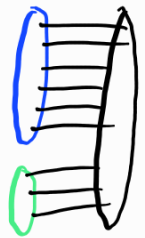
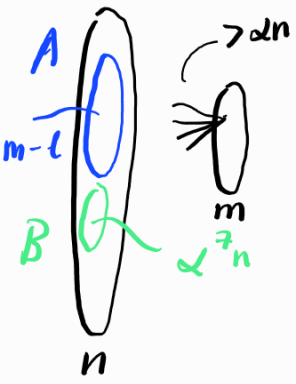
Absorbing: from almost finished to finished.

$\triangle \triangle \triangle \dots \triangle$ n colours

Key lemma

$m = \Omega^2 n$ $\ell = \Omega^9 n$

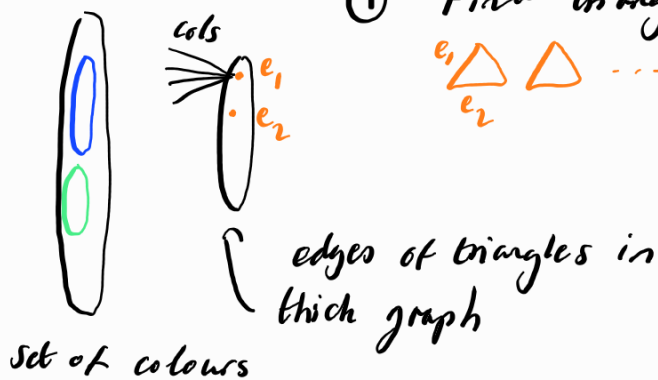
For any $B_0 \subseteq B$ of size ℓ there is a perfect matching from $A \cup B_0$.



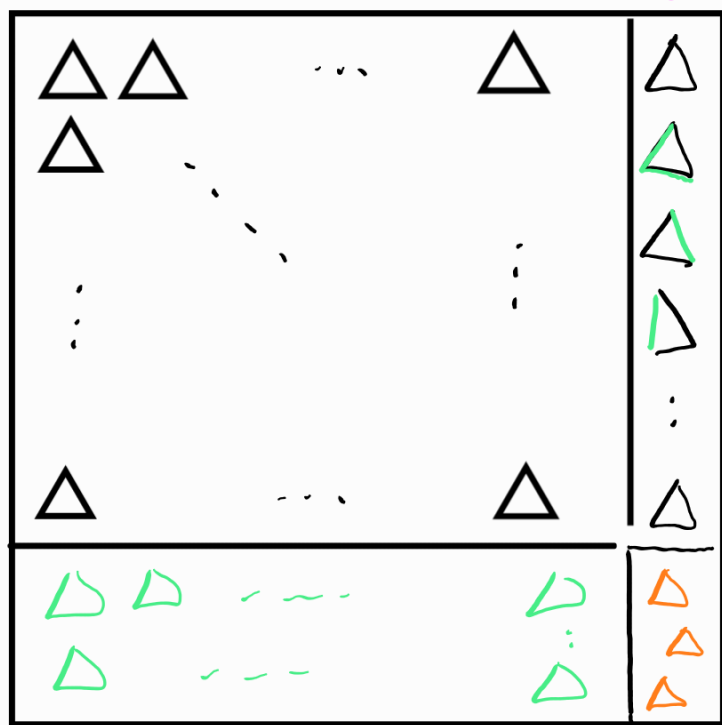
Find **absorber**.

Many edges in **thick graph** (edges of many colours).

① Find triangles in thick graph.



② Use most black colours



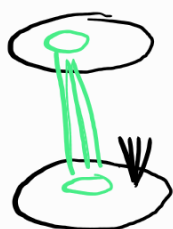
③ Use remaining black + green

④ Use remaining vertices (green cols)

⑤ Use the lemma to assign remaining colours

Blow-up lemma (Komlós - Sárközy - Szemerédi)

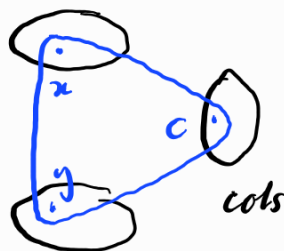
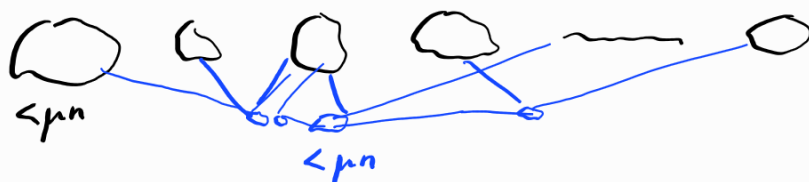
(+ regularity lemma) is a very useful tool for embedding spanning graphs



'quasirandom': every not-too-small pair spans roughly the same proportion of edges as globally
reasonable min degree

→ contains any bounded degree graph 'that fits'

(Cheng - S.) Transversal blow-up lemma for ^{bounded degree} separable graphs



'quasirandom'
reasonable min degree

→ contains any bounded degree separable graph 'that fits'

Use TBUL in place of BUL to get transversal versions of embedding results.

Open questions

($\binom{r}{2}$) graphs G_i , $e(G_i) > ? \Rightarrow$ rainbow K_r ?

Spanning: exact results $(n + \epsilon n)$

Other colour patterns between monochromatic and rainbow