# Permutations that separate close elements 

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## Torus packings



## Algebraic phrasing

For $i, j \in \mathbb{Z}_{n}$, let $\|i, j\|_{n}$ be the distance between $i$ and $j$ when the elements of $\mathbb{Z}_{n}$ are written in a circle.

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## Definition (An overlapping rectangle)

A permutation $\pi: \mathbb{Z}_{n} \rightarrow \mathbb{Z}_{n}$ has an $(s, k)$-clash if there exist distinct $i, j \in \mathbb{Z}_{n}$ with $\|i, j\|_{n}<s$ and $\|\pi(i), \pi(j)\|_{n}<k$.

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## Definition (No overlapping rectangles)

A permutation $\pi: \mathbb{Z}_{n} \rightarrow \mathbb{Z}_{n}$ is $(s, k)$-clash-free if it has no $(s, k)$-clashes.

## Related work

- Generalisations of $k=2$ case: cyclic matching sequencability for graphs: Alspach, Bull. ICA 2008, Brualdi-Kiernan-Meyer, Australas. J. Comb. 2012; Kreher-Pastine-Tollefson, Australas. J. Comb. 2015.


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- Non-cyclic case (cylinder or square, not torus): Mammoliti-Simpson, Australas. J. Comb. 2020.


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- Non-cyclic case (cylinder or square, not torus): Mammoliti-Simpson, Australas. J. Comb. 2020.
- Packing diamonds rather than rectangles (large distance in the Manhattan metric): Aspvell-Liang Stanford Tech. Report 1980; Bevan-Homberger-Tenner JCT-A 2018; SRB-Homberger-Winkler JCT-A 2019.


## The main question

Definition (How wide can rectangles be?)
Let $n$ and $k$ be fixed. Define $\sigma(n, k)$ to be the largest $s$ such that an $(s, k)$-clash-free permutation $\pi$ of $\mathbb{Z}_{n}$ exists.

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We can't have sk $=n$ :


So $s k \leq n-1$.

## Mammoliti-Simpson conjecture

Theorem (SRB, JCT-A 2023)

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\lfloor(n-1) / k\rfloor-1 \leq \sigma(n, k) \leq\lfloor(n-1) / k\rfloor
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Set $\rho(0)=0, \rho(1)=12$, and so on. $\rho$ is $(k, s)$-clash-free. Set $\pi=\rho^{-1}$. Then $\pi$ is ( $s, k$ )-clash-free.

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## Theorem (SRB-Etzion, 2023+)

Let $n$ and $k$ be fixed positive integers, with $k<n$. Write $s=\lfloor(n-1) / k\rfloor$, so $n=s k+r$ where $1 \leq r \leq k$.

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- If $r \geq s$ or $k=r$, then $\sigma(n, k)=\lfloor(n-1) / k\rfloor$.
- If $r<s$ and $r<k$ and $d_{s} d_{k}$ divides $n$, then $\sigma(n, k)=\lfloor(n-1) / k\rfloor$.


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- If $r \geq s$ or $k=r$, then $\sigma(n, k)=\lfloor(n-1) / k\rfloor$.
- If $r<s$ and $r<k$ and $d_{s} d_{k}$ divides $n$, then $\sigma(n, k)=\lfloor(n-1) / k\rfloor$.
- If $r<s$ and $r<k$ and $d_{s} d_{k}$ does not divide $n$, then $\sigma(n, k)=\lfloor(n-1) / k\rfloor-1$.


## A sketch proof

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In every row, and every column, exactly $r$ positions are uncovered.
Every rectangle touches 4 others, one on each side:


Rectangles form east-west and north-south lines: warp and weft threads. Threads cannot change direction:


## A sketch proof 2

Threads must be periodic, giving the condition that $d_{s} d_{k}$ divides $n$.

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Can classify permutations by jumpers: two sequences determining sizes of gaps.

## Jumpers

## Definition

An $(s, k, n)$-jumper is a pair $\left(\left(a_{i}\right),\left(b_{i}\right)\right)$ of sequences of integers with the following properties:
(1) ( $a_{i}$ ) has period dividing $d_{s}$, and $\left(b_{i}\right)$ has period dividing $d_{k}$.
(2) We have $1 \leq a_{i}<s$ and $1 \leq b_{i}<k$ for $i \geq 0$.
(3) The $d_{k}$ partial sums $\sum_{i=0}^{\ell-1} b_{i}$ where $0 \leq \ell<d_{s}$ are distinct modulo $d_{k}$. Moreover, $d_{s} d_{k}$ divides $\sigma_{b}$ where $\sigma_{b}=\sum_{i=0}^{d_{k}-1} b_{i}$.
(1) The $d_{s}$ partial sums $\sum_{i=0}^{m-1} a_{i}$ where $0 \leq m<d_{s}$ are distinct modulo $d_{s}$. Moreover, $d_{s} d_{k}$ divides $\sigma_{a}$ where $\sigma_{a}=\sum_{i=0}^{d_{s}-1} a_{i}$.
(0) Defining $\sigma_{a}$ and $\sigma_{b}$ as above, $\sigma_{a} \sigma_{b}=d_{s} d_{k} r$.

## The classification

## Theorem

Let $n$ and $k$ be fixed integers with $k<n$. Set $s=\lfloor(n-1) / k\rfloor$, and define $r$ by $n=s k+r$ for $1 \leq r \leq k$. Define $d_{s}=\operatorname{gcd}(n, s)$ and $d_{k}=\operatorname{gcd}(n, k)$. Assume that $r<k$ and $r<s$. Futhermore, suppose that $d_{s} d_{k}$ divides $n$.

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## Thanks!

